

COURSE NAME: ALGORITHMS

COURSE CODE: CIS 212

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LINEAR SEARCH

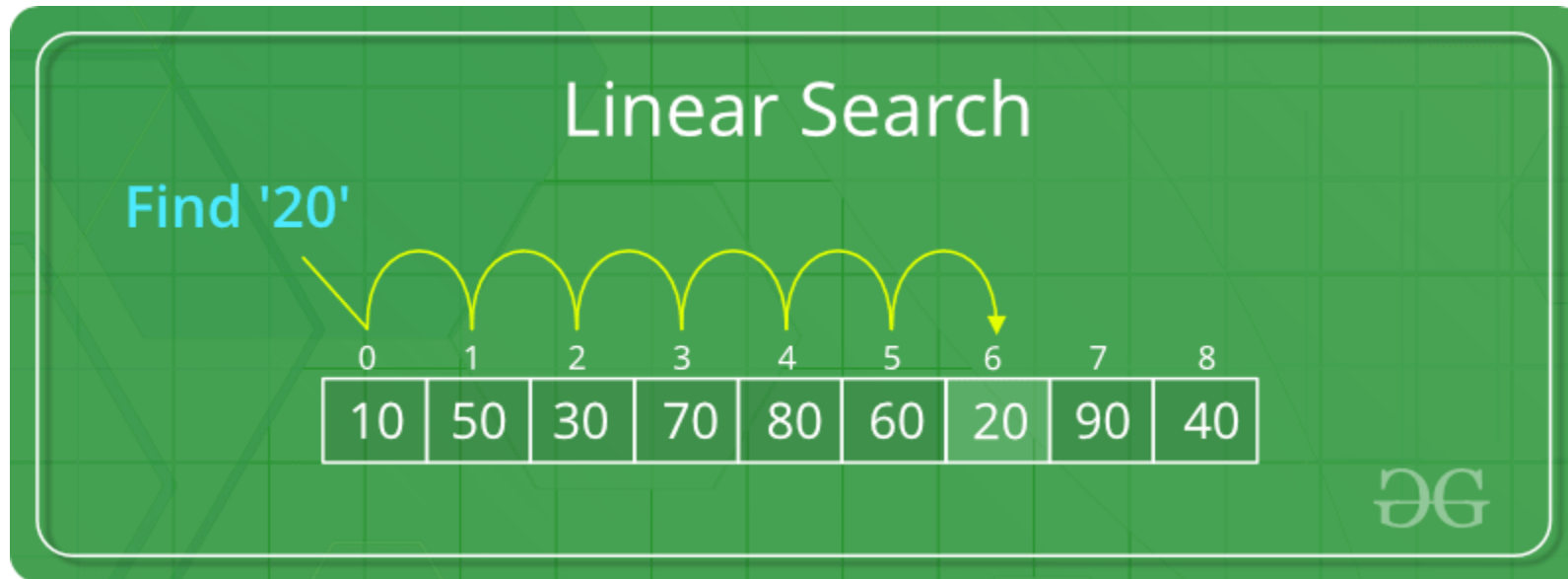
- Linear Search or Sequential Search.
- Worst Case TC: $O(n)$
- Best Case TC: $O(1)$
- Average Case TC: $(n+1)/2$
- Explanation of Average Case: $1+2+3+\dots+n = n*(n+1)/2 = (n+1)/2$
- Worst Case TC in Iterative: $O(1)$ (when elements in the first or last position)
- Space Complexity: $O(1)$

Advantages & Disadvantages

- **Advantages:**
 - Perform faster in small and medium size array
 - The array does not need to be sorted
 - Not affected by insertions and deletions.
- **Disadvantages:**
 - Slow searching in large arrays

LINEAR SEARCH

- Start from the leftmost element of arr[] and one by one compare x with each element of arr[]
- If x matches with an element, return the index.
- If x doesn't match with any of elements, return -1



BINARY SEARCH

- Binary Search, half-interval search, logarithmic search, binary chop
- Worst Case TC: $O(\log n)$
- Best Case TC: $O(1)$
- Average Case TC: $(\log n)$
- Space Complexity: $O(1)$

BINARY SEARCH

- **Worst Case TC explanation:**
- **Iteration-1:** length of array= n
- **Iteration-2:** length of array= $n/2$
- **Iteration-3:** length of array= $n/2^2$
- **Iteration-k:** length of array= $n/2^k$
- **After k iterations**, the length of array becomes 1, so
- **Length of array:** $n/2^k=1$
- $n=2^k$
- $\log_2(n)=\log_2(2^k)$
- $\log_2(n)=k \log_2(2) // \log_a(a)=1$
- **$K=\log_2(n)$**

BINARY SEARCH

- Compare x with the middle element.
- If x matches with the middle element, we return the mid index.
- Else If x is greater than the mid element, then x can only lie in the right half subarray after the mid element. So we recur for the right half.
- Else (x is smaller) recur for the left half.

